

**Discrete Optimization** (Spring 2018)

**Assignment 7**

**Problem 3** can be **submitted** until April 20 12:00 noon into the box in front of MA C1 563.  
You are allowed to submit your solutions in groups of at most three students.

**Problem 1**

Determine the dual program for the following linear program:

$$\begin{aligned} \min \quad & 3x_1 + 2x_2 - 3x_3 + 4x_4 \\ & 2x_1 - 2x_2 + 3x_3 + 4x_4 \leq 3 \\ & \quad \quad \quad x_2 + 3x_3 + 4x_4 \geq -5 \\ & 2x_1 - 3x_2 - 7x_3 - 4x_4 = 2 \\ & \quad \quad \quad x_1 \geq 0 \\ & \quad \quad \quad x_4 \leq 0 \end{aligned}$$

**Solution:**

After some simplifications, one obtains:

$$\begin{aligned} \max \quad & 3y_1 - 5y_2 + 2y_3 \\ & 2y_1 + 2y_3 \leq 3 \\ & -2y_1 + y_2 - 3y_3 = 2 \\ & 3y_1 + 3y_2 - 7y_3 = -3 \\ & \quad \quad \quad y_1 + y_2 - y_3 \geq 1 \\ & \quad \quad \quad y_1 \leq 0 \\ & \quad \quad \quad y_2 \geq 0 \end{aligned}$$

**Problem 2**

In the GitHub repository of the course, in the ‘Programming’ folder, you will find the file ‘Simplex.py’ which contains the code to run the simplex algorithm on non-degenerate linear programs in standard form. However, some code is missing. Complete the code by filling the gaps (indicated by underscores) and send it to igor@malinovic.epfl.ch.

**Solution:**

The solution can be found in the GitHub repository of the course.

**Problem 3** (★)

Suppose you are given an oracle algorithm, which for a given polyhedron

$$P = \{\tilde{x} \in \mathbb{R}^{\tilde{n}} : \tilde{A}\tilde{x} \leq \tilde{b}\}$$

gives you a feasible solution or asserts that there is none.

Consider the LP  $\max\{c^T x : Ax \leq b, x \in \mathbb{R}^n\}$ , and assume it is feasible and bounded. Show that one can obtain an optimal solution of the LP using a single call of the oracle algorithm on a suitable polyhedron. *Hint: Use duality theory!*

**Solution:**

The LP is feasible and bounded, thus an optimum solution must exist. Strong duality tells us that

the dual  $\min\{b^T y : A^T y = c, y \geq 0\}$  is feasible and bounded. For optimal solutions  $x^*$  of the primal and  $y^*$  of the dual we have  $b^T y^* = c^T x^*$ .

Thus every point  $(x^*, y^*)$  of the polyhedron

$$\begin{aligned} c^T x &= b^T y \\ Ax &\leq b \\ A^T y &= c \\ y &\geq 0 \end{aligned}$$

is optimal. Hence with one oracle call for the polyhedron above we get an optimal solution of the LP.

#### Problem 4

Consider the following linear program:

$$\begin{aligned} \max \quad & x_1 + x_2 \\ \text{subject to} \quad & 2x_1 + x_2 \leq 6 \\ & x_1 + 2x_2 \leq 8 \\ & 3x_1 + 4x_2 \leq 22 \\ & x_1 + 5x_2 \leq 23 \end{aligned}$$

Show that  $(4/3, 10/3)$  is an optimal solution by using weak duality.

#### Solution:

The assignment  $(4/3, 10/3)$  has the objective function value of  $14/3$ . In order to prove that it is optimal (via strong duality), we are going to form the dual LP, and find a feasible solution to the dual that achieves the same objective value. The dual is:

$$\begin{aligned} \min \quad & 6y_1 + 8y_2 + 22y_3 + 23y_4 \\ \text{subject to} \quad & 2y_1 + y_2 + 3y_3 + y_4 = 1 \quad (1) \\ & y_1 + 2y_2 + 4y_3 + 5y_4 = 1 \quad (2) \\ & y_1, y_2, y_3, y_4 \geq 0 \end{aligned}$$

Thus, we are looking for a feasible dual solution such that  $6y_1 + 8y_2 + 22y_3 + 23y_4 = 14/3$ . By using Gaussian elimination on this constraint combined with (1) and (2) we get:

$$\begin{aligned} -4y_2 - 2y_3 - 7y_4 &= -4/3, \\ -3y_2 - 5y_3 - 9y_4 &= -1 \end{aligned}$$

and further

$$14/3y_3 + 5y_4 = 0.$$

Since  $y_1, y_2, y_3, y_4 \geq 0$ , we have that  $y_3 = y_4 = 0$  and then  $y_1 = y_2 = 1/3$ . This is the desired feasible dual solution coinciding with the primal solution  $(4/3, 10/3)$ , proving the optimality of the latter.

#### Problem 5

Let  $P = \{x \in \mathbb{R}^n : Ax \leq b\}$  a bounded, non-empty polyhedron. Formulate a linear program that computes the largest ball inside  $P$ .

**Solution:**

A ball of radius  $r$  and center  $x$  is contained in  $P$  if and only if  $x \in P$  and  $x$  has distance at least  $r$  from any hyperplane defining  $P$ . Hence we obtain the following linear program:

$$\begin{aligned} \max \quad & r \\ \text{subject to} \quad & \frac{b_i - a_i x}{\|a_i\|} \geq r \quad \forall i = 1, \dots, m \\ & Ax \leq b \end{aligned}$$

where  $a_1, \dots, a_m$  are the rows of  $A$  and  $b = (b_1 \dots b_m)^\top$ .