

**Discrete Optimization** (Spring 2017)

**Assignment 7**

**Problem 3** can be **submitted** until April 14 12:00 noon into the right box in front of MA C1 563.

You are allowed to submit your solutions in groups of at most three students.

**Problem 1**

Provide a proof or a counterexample to the following statements:

- i) An iteration of the simplex method cannot move the feasible solution by a strictly positive distance while leaving the objective function value unchanged.
- ii) If  $B$  is an optimal basis, then all the components of  $\lambda_B$  are strictly positive.

**Problem 2**

Let  $B$  be an optimal basis. Provide a proof or a counterexample to the following statements:

- i) If  $\lambda_B$  is strictly positive then the optimal solution is unique.
- ii) If the optimal solution is unique then  $\lambda_B$  is strictly positive.

**Problem 3** (★)

Let  $\max\{c^T x : x \in \mathbb{R}^n, Ax \leq b\}$  be a non-degenerate linear program with a unique optimal solution. Prove that  $\lambda_B$  is strictly positive for the corresponding optimal basis  $B$ .

**Problem 4**

Recall that a bounded polyhedron  $P \subseteq \mathbb{R}^n$  is called a *polytope*, i.e. there exists an  $M \in \mathbb{R}$  such that  $P \subseteq [-M, M]^n$ . In the following let  $P$  be a non-empty polytope.

- (i) Prove that  $P$  has vertices.
- (ii) Let  $u_1, \dots, u_\ell \in \mathbb{R}^n$  be the vertices of  $P$ . Show that  $P$  is equal to the convex hull  $\text{conv}(\{u_1, \dots, u_\ell\})$ , without using the separating hyperplane theorem.

*Hint:* To show that  $P \subseteq \text{conv}(\{u_1, \dots, u_\ell\})$  one possibility is to suppose that there exists a point  $x^* \in P \setminus \text{conv}(\{u_1, \dots, u_\ell\})$  and consider the following linear program and its dual:

$$\begin{array}{ll} \min & 0^T \cdot \lambda \\ \text{s.t.} & \sum_{i=1}^{\ell} \lambda_i u_i = x^* \\ & \sum_{i=1}^{\ell} \lambda_i = 1 \\ & \lambda \geq 0 \\ & \lambda \in \mathbb{R}^{\ell} \end{array} \qquad \begin{array}{ll} \max & (x^*)^T \cdot c + \beta \\ \text{s.t.} & u_i^T \cdot c + \beta \leq 0 \quad i = 1, \dots, \ell \\ & c \in \mathbb{R}^n, \beta \in \mathbb{R} \end{array}$$

Conclude that the primal problem (the minimization) must be infeasible and that its dual problem

(the maximization) is unbounded. From this, yield a contradiction to the fact that there always exists an optimal vertex in  $P$  for any linear optimization problem over  $P$ .

**Problem 5**

Given two polynomials  $p_1(x) = \sum_{i=0}^n a_i x^i$ ,  $p_2(x) = \sum_{i=0}^m b_i x^i$ , we say that  $p_1$  is lexicographically smaller than  $p_2$ , and we write  $p_1 <_L p_2$ , if  $(a_0, a_1, \dots, a_n) <_L (b_0, b_1, \dots, b_m)$  (i.e. if  $i$  is the first index such that  $a_i$  is different than  $b_i$ ,  $a_i < b_i$ ). Prove that for any  $p_1, p_2$  there is an  $\epsilon^* > 0$  such that the following holds:

$$p_1(\epsilon) < p_2(\epsilon) \quad \forall \epsilon : 0 < \epsilon < \epsilon^* \quad \Leftrightarrow \quad p_1 <_L p_2.$$