

# Convexity

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## Assignment Sheet 6 - Solutions

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### Exercise 1

Let  $P$  be a polyhedron and  $F$  a facet. Show  $\dim(F) = \dim(P) - 1$ .

#### Solution:

Let  $P$  be described by the non-redundant system  $Ax \leq b$  and  $F$  be a facet of  $P$ . Let  $A^=x \leq b^=$  the system of implicit equalities,  $a^T x \leq \beta$  the inequality corresponding to  $F$  (by the theorem of the lecture), and  $A'x \leq b'$  be the rest of the system.

In the proof of the theorem linking the facets to the inequalities, we saw that there is a point  $x^* \in P$  with

$$\begin{pmatrix} A^= \\ a^T \\ A' \end{pmatrix} x = \begin{pmatrix} b^= \\ \beta \\ b' \end{pmatrix}.$$

As  $F = \{x \in P : a^T x = \beta\}$  and all inequalities in  $A'x \leq b'$  are no implicit equalities, the system of implicit equalities of  $F$  is exactly one larger than of  $P$ . Once we have shown that the rank also grew by one, we are done using the characterization of the dimension of a polyhedron. But if the vector  $a$  was linearly dependent of the rows in  $A^=$ , the equality  $a^T x = \beta$  would be fulfilled by all elements in  $P$ , hence  $F$  would not be a facet.

### Exercise 2

Prove the Birkhoff- von Neumann Theorem:

The extreme points of the set of doubly stochastic matrices

$$M = \left\{ A = (a_{i,j})_{1 \leq i,j \leq n} \in \mathbb{R}_{\geq 0}^{n \times n} \mid \sum_{i=1}^n a_{i,j_0} = \sum_{j=1}^n a_{i_0,j} = 1 \forall 1 \leq i_0, j_0 \leq n \right\}$$

are precisely the permutation matrices, i.e.  $M \cap \{0, 1\}^{n \times n}$ .

#### Solution:

Note that  $M$  is a polytope with constraints

$$\sum_{i=1}^n a_{ij} = 1 \quad \text{for } j = 1, \dots, n \quad (1)$$

$$\sum_{j=1}^n a_{ij} = 1 \quad \text{for } i = 1, \dots, n$$

$$a_{ij} \geq 0 \quad \text{for } i, j = 1, \dots, n \quad (2)$$

for the contained matrices  $A = (a_{ij})_{i,j}$ . Each permutation matrix is extreme, as otherwise at least one 1-entry would be a convex combination of some numbers in  $[0, 1)$ , which is not possible. If we can show that each extreme point of  $P$  is a convex combination of permutation matrices, we are basically done.

We will proceed inductively. For  $n = 1$ , the scalar 1 is the only doubly-stochastic matrix, and a permutation matrix. Let  $A$  be a vertex of the polytope described by (1). Hence,  $A$  has to fulfil  $n^2$  linearly independent inequalities with equality. (Recall the characterization of extreme points.) As the first  $2n$  inequalities are linearly independent, at least  $n^2 - 2n + 1$  of the other inequalities have to be fulfilled with equality, implying that at least  $n^2 - 2n + 1$  of the entries in  $A$  are zero. By the pigeon hole principle, at least one row of  $A$  has to consist of  $n - 1$  zeroes and one entry equal to 1, say w.l.o.g.  $a_{11} = 1$ . But then the first column contains  $n - 1$  zeroes as well and the matrix  $A'$  arising from  $A$  by deleting the first row and column is doubly stochastic and of smaller dimension, hence it is the convex combination of permutation matrices.

### Exercise 3 [★]

Using the separation theorem, prove that the system  $Ax = b, x \geq 0$  has no solution if and only if there is a vector  $c$  s.th.  $c^T A \geq 0$  and  $c^T b < 0$ .

#### Solution:

Assume the system has no solution. All  $b'$  for which there exists an  $x \geq 0$  with  $Ax = b'$  are in the cone generated by the columns of  $A$ , call this cone  $S$ . As  $b \notin S$ , there exists a separating hyperplane  $\{x : c^T x = \gamma\}$  with  $c^T b < \gamma$  and  $c^T y \geq \gamma$  for any  $y \in S$ . As  $S$  is a cone, we can replace  $\gamma$  by 0 (in particular,  $\gamma \leq 0$ ). Hence,  $c$  is the desired vector.

On the other hand, if there exists such a vector it can be interpreted as a separating hyperplane. Or, to give it another angle, if  $c^T x = 0$  is the hyperplane and there was a solution  $x^*$ ,  $0 > c^T b = c^T A x^* \geq 0$  yields a contradiction.

### Exercise 4

Let  $F$  be an (inclusion-wise) minimal face of a polyhedron  $P = \{x : Ax \leq b\}$ . Show that the following holds.

$$\forall x, y \in F : Ax = Ay.$$

#### Solution:

Let  $A'x \leq b'$  be a maximum subsystem s.th. for all  $x, y \in F$  we have  $A'x = A'y$  and let  $A''x \leq b''$  be the remaining system. If the remaining system is empty, we are done.

Otherwise, we want to find a proper face of  $F$ . By transitivity, this will be a face of  $P$ , contradicting minimality.

Pick  $x, y \in F$  s.th. there is at least one inequality with  $a^T x \neq a^T y$  in the system  $A''x \leq b''$ . Consider  $z(\lambda) := x + \lambda(y - x)$ . First note that  $A'z(\lambda) = A'x$  for any  $\lambda$ . As  $z(0)$  is in  $F$  and due to the choice of  $x, y$ , there is a  $\lambda_0$  with  $z(\lambda)$  being feasible for  $A''x \leq b''$  and one of the inequalities is fulfilled with equality, say  $a_1^T z(\lambda_0) = \beta_1$ . (Metaphorically speaking, we start at  $x$  and, by altering  $\lambda$ , walk in the direction of  $y$  or straight away from  $y$  for negative  $\lambda$  until we hit the boundary of  $F$ .) But then  $F \cap \{x : a_1^T x = \beta_1\}$  is a non-empty intersection of  $F$  with a supporting hyperplane, yielding the contradiction.